

On Clustering Induced Voronoi Diagrams^{*}

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1 Overview

Voronoi diagram is a fundamental geometric structure with numerous applications in many different areas [1–3]. Ordinary Voronoi diagram is a partition of the space \mathbb{R}^d into a set of cells induced by a set P of points (or other types of objects) called sites, where each cell c_i of the diagram is the union of all points in \mathbb{R}^d which have a closer (or farther) distance to a site $p_i \in P$ than to any other sites. In some sense, cells in a Voronoi diagram can be viewed as formed by competitions among all sites in \mathbb{R}^d such that the winner site for any point $q \in \mathbb{R}^d$ is the one having a larger “influence” on q defined by its distance to q .

In this paper, we generalize the concept of Voronoi diagram to *Clustering Induced Voronoi Diagram* (CIVD). In CIVD, we consider a set P of n points (or other types of objects) and a non-negative influence function F which measures the joint influence $F(C, q)$ from each subset C of P to any point q in \mathbb{R}^d . The Voronoi cell of C is the union of all points in \mathbb{R}^d which receive a larger influence from C than from any other subset $C' \subseteq P$. This means that CIVD considers all subsets in the power set $U = 2^P$ of P as its sites (called *cluster sites*), and partitions \mathbb{R}^d according to their influences. While CIVD in general can have exponentially many cells, it is possible that for some interesting influence functions only a small number of subsets in U have non-empty Voronoi cells, making the problem solvable. As application of our model and technique, we consider two representative CIVD problems, vector CIVD and density-based CIVD.

Relation to Previous Works: To our best knowledge, there is no previous work on the general CIVD problem. Our CIVD model obviously extends the ordinary Voronoi diagrams [2], where each site is a one-point cluster. (Note that the ordinary Voronoi diagrams can be viewed as special CIVDs equipped with proper influence functions.) Some Voronoi diagrams [3, 12] allow a site to contain multiple points, such as the k -th order Voronoi diagram [3]. Some two-point site Voronoi diagrams were also studied [4, 5, 7, 8, 10, 11, 13], in which each site has exactly two points. Obviously, such Voronoi diagrams are different from CIVD.

For Vector CIVD, influence between any two points p and q is defined by a force-like vector. The problem is related to the N-body problem [9], which shares with the Vector CIVD problem a similar idea of modeling joint force by influence functions. Density-based CIVD enables us to generate all density-based clusters as well as their approximate Voronoi cells. The problem is related to density-based clustering which is widely used in many applications.

1.1 Results and Techniques

The main result of the paper is a general technique called Approximate Influence(AI) Decomposition, which can be used to generate $(1 - \epsilon)$ -approximate CIVD. We also apply AI decomposition to develop assignment algorithms for vector CIVD and density base CIVD. Below is a list of our main results.

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- **Properties of Influence Function:** We investigate the general and sufficient conditions which allow the influence function to yield only a small number of non-empty approximate Voronoi cells. We show that the following three properties are sufficient: **Similarity Invariant property, Locality property, Local Domination property.** The first property means that for any point $q \in \mathbb{R}^d$, its maximum influence cluster site remains the same after a similarity transformation about it. The second property indicates that a small perturbation on a cluster site C or q only changes slightly their influence. The third property implies that each cluster site may have dominating influence in its neighborhood.
- **Approximate Influence(AI) Decomposition:** We present a standalone technique called *approximate influence decomposition* (or AI decomposition) for general CIVD problems. In $O(n \log n)$ time, this technique partitions the space \mathbb{R}^d into a nearly linear number (*i.e.*, $O(n \log n)$) of cells so that for each such cell c , there exists a (possibly unknown) subset $C \subseteq P$ whose influence to any point $q \in c$ is within a $(1 - \epsilon)$ -approximation of the maximum influence that q can receive from any subset of P , where $\epsilon > 0$ is a fixed small constant. In this technique, we also develop a new data structure called *box-clustering tree*, based on an extended quad-tree decomposition and guided by a *distance-tree* built from the well-separated pair decomposition [6]. In some sense, our AI decomposition may be viewed as a generalization of the well-separated pair decomposition.
- **Assignment Algorithms for Vector CIVD and Density-base CIVD:** Based on the AI Decomposition, we develop assignment algorithms for the Vector CIVD and the Density-based CIVD problems. Particularly, we show that it is possible to determine a proper cluster site for each cell in the decomposition and form a $(1 - \epsilon)$ -approximate CIVD for each problem for any given small constant $\epsilon > 0$. For Vector CIVD, the assignment algorithm is based on several new techniques such as aggregation-tree and majority path decomposition, and runs in $O(n \log^{d+1} n)$ time. For Density-based CIVD, the assignment algorithm takes $O(n \log^2 n)$ time and can be obtained from a modification of the AI decomposition.

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